

5CCS2FC2: Foundations of Computing II

P vs NP: The Million Dollar Problem

Week 3

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The Cook-Levin Theorem

(SAT is NP-complete)

The Cook-Levin Theorem

Theorem (The Cook-Levin Theorem) The Boolean Satisfiability problem **SAT** is **NP**-complete.

Proof: We want to show that **every** problem $X \in \mathbf{NP}$ can be reduced to **SAT**.

Step 1) Let $X \in \mathbf{NP}$ be **any problem** belonging to the class **NP**.

By **definition** there is some NDTM M such that

M has a polynomial-length
computation that accepts w $\iff w \in X$

The Cook-Levin Theorem

Step 2) For each **state** $q \in Q$, **symbol** $a \in \Sigma$ and **integers** $i, t \leq T_M(n)$

$C_{i,t,a}$ = Cell i of the M 's tape contains symbol a at time t

$H_{i,t}$ = The tape head is in position i at time t

$S_{q,t}$ = The machine is in state q at time t

The Cook-Levin Theorem

- “The machine cannot be both in state q_4 and state q_5 at time $t = 2$ ”

$$\neg(S_{q_4,2} \wedge S_{q_5,2})$$

- “At time $t = 3$ the tape head must be in positions 0 or 1 or 2 or 3”

$$(H_{0,3} \vee H_{1,3} \vee H_{2,3} \vee H_{3,3})$$

The Cook-Levin Theorem

- "At time $t = 3$ if cell 1 contains an a then it must not contain a b "

$$C_{3,1,a} \rightarrow \neg C_{3,1,b}$$

- "At time $t = 6$ if the machine is in state q_1 and the tape head is in position 4 and cell 4 contains an a at time $t = 6$ then at time $t = 7$ the machine will be in state q_2 , the tape head will move to position 3 and cell 4 will contain a b ."

$$(S_{q_1,6} \wedge H_{4,6} \wedge C_{4,6,a}) \rightarrow (S_{q_2,7} \wedge H_{3,7} \wedge C_{4,7,b})$$

The Cook-Levin Theorem

Step 3) With these **propositional variables**, we can construct a formula:

$$F_{M,w} = \begin{array}{l} M \text{ has a polynomial-length} \\ \text{computation that accepts } w \end{array}$$

$$\begin{array}{l} M \text{ has a polynomial-length} \\ \text{computation that accepts } w \end{array} \iff w \in X$$

The Cook-Levin Theorem

Step 4) We then have that

$$w \in X \iff F_{M,w} \text{ is satisfiable}$$

$$X \leq_p \text{SAT}$$

$$X \leq_p \text{SAT} \quad \text{for all } X \in \text{NP}$$

Q.E.D.

The Cook-Levin Theorem

Theorem **SAT** is **NP**-complete.

End of Slides!

