

5CCS2FC2: Foundations of Computing II

# Introduction to Graph Algorithms

Week 4

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# Graph Algorithms

## Graph Algorithms

- Graph Algorithms

- We expand the notation for the **termination time** for algorithms on graphs

$$T(|V|, |E|) = \text{The number of } \textit{primitive operations} \text{ on a graph with } |V| \text{ vertices and } |E| \text{ edges}$$

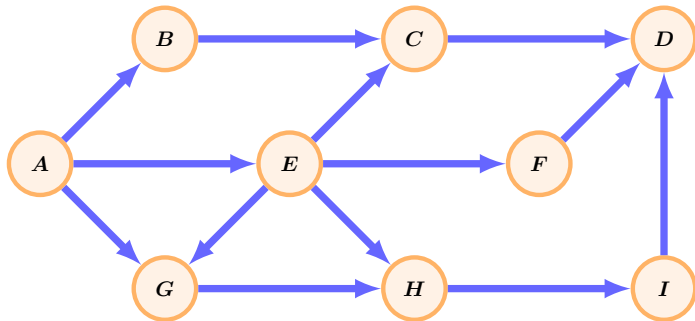
- Our **primitive operations** are operations we can perform in  $O(1)$  steps.

## Breadth-First-Search

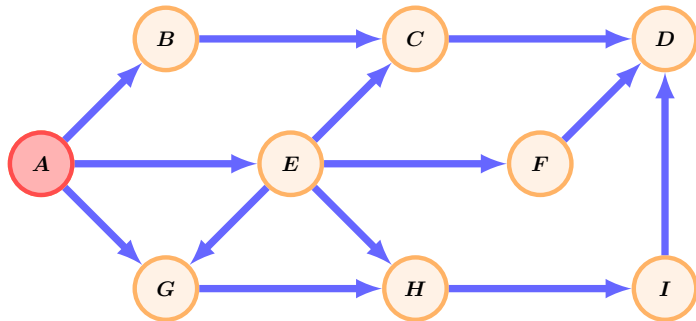
### BREATH-FIRST-SEARCH :

1. Select root vertex and add to a queue
2. **While** queue is not-empty:
  3. Inspect the first element  $u$  on the queue.
  4. Add any unvisited neighbours of  $u$  to the queue.
  5. Remove  $u$  from the queue.

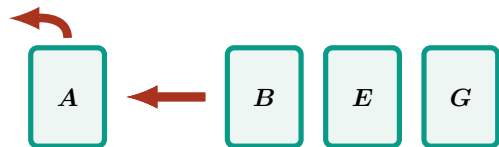
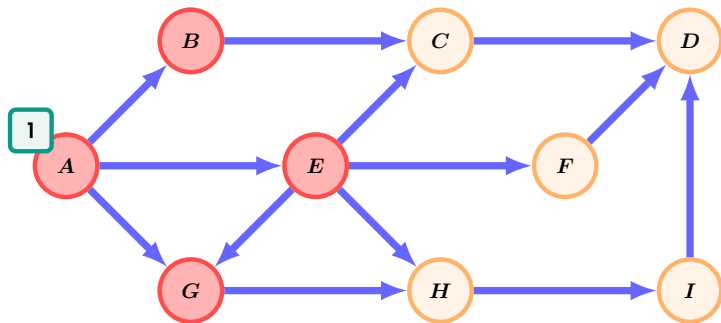
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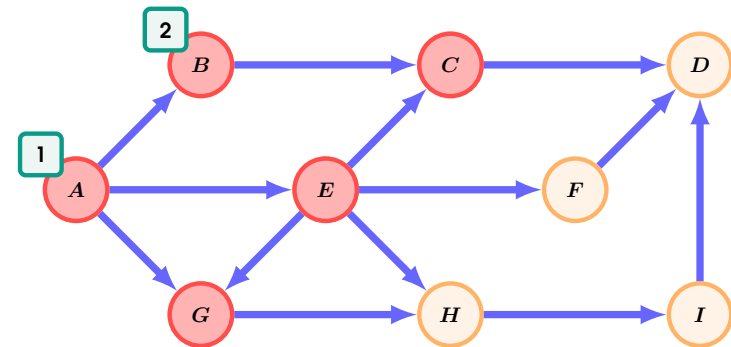
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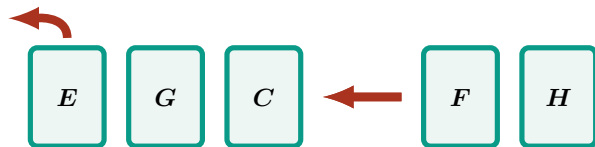
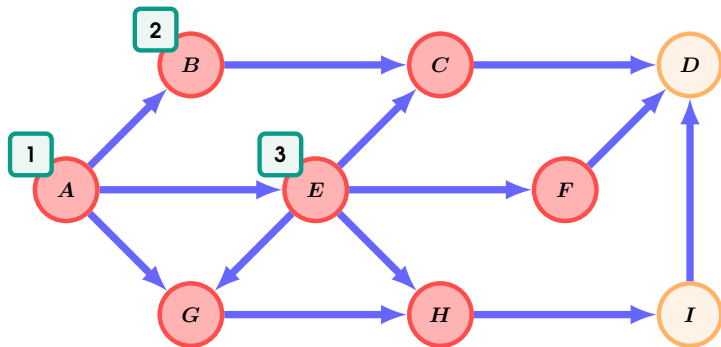
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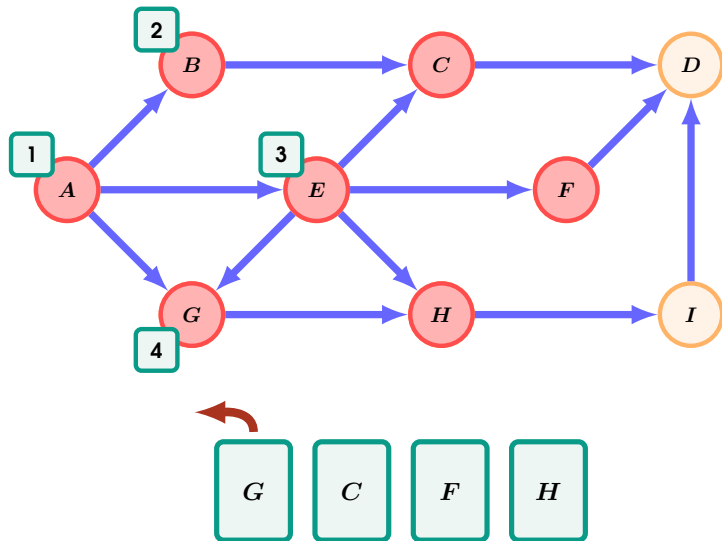
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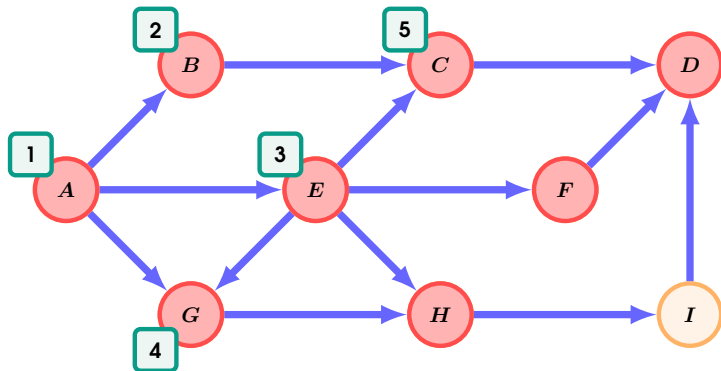
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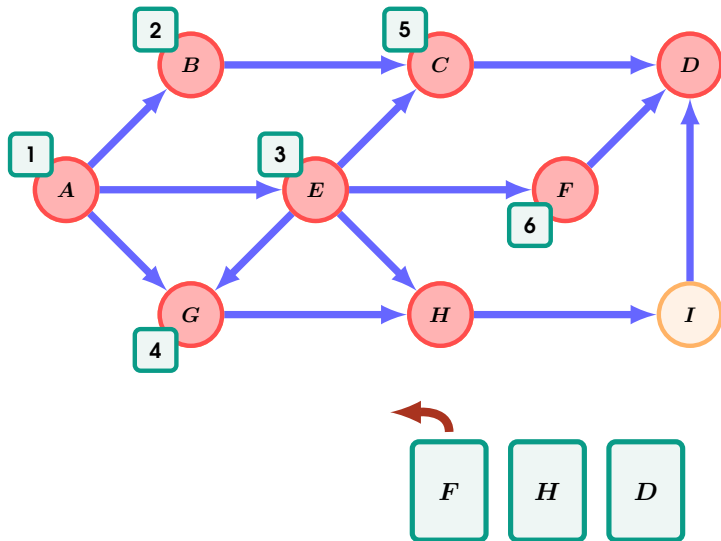
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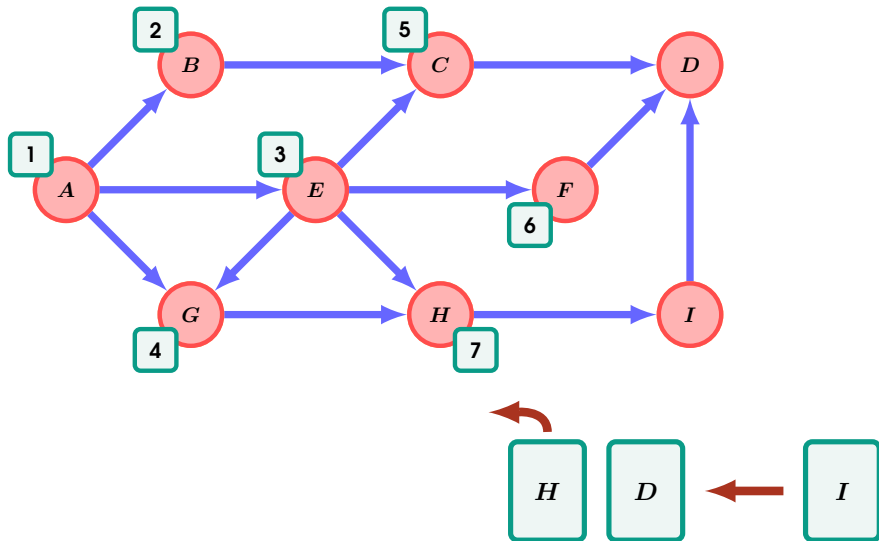
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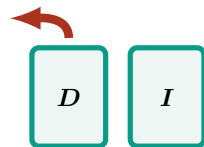
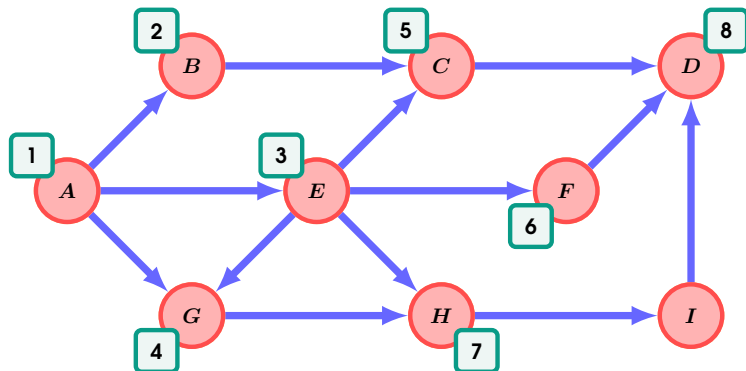
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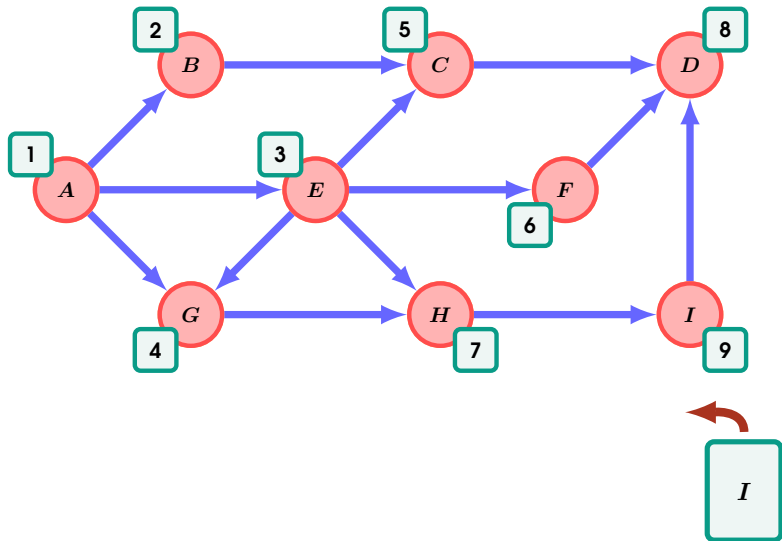
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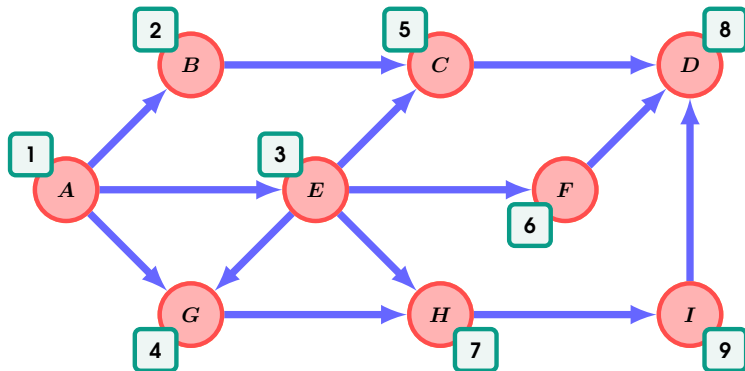
## Breadth-First-Search



## Breadth-First-Search



## Breadth-First-Search



**Queue Empty so Terminate!**

## Breadth-First-Search

**Theorem** The termination time for Breadth-First-Search is  $O(|V| + |E|)$ .

### *Proof:*

- Note that every vertex is **enqueued** and **dequeued** at most once.  
So **Steps 3–5** are repeated once for each  $u \in V$ .
- Inspecting (**Step 3**) and Adding (**Step 5**) each take at most  $O(1)$  operations.
- However, **Step 4** requires  $O(|\text{Adj}(u)|)$  operations to queue each of the adjacent vertices to  $u$ .

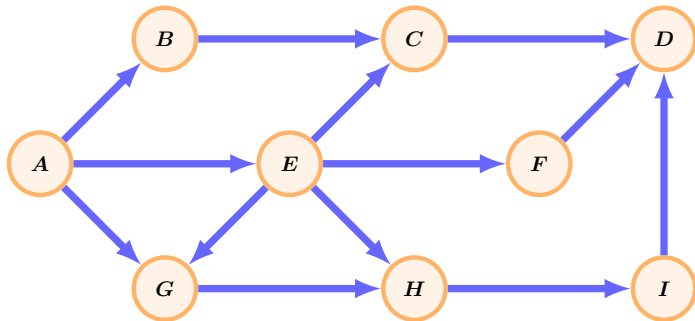
(where  $\text{Adj}(u)$  denotes the set of vertices adjacent to  $u$ )

## Depth-First-Search

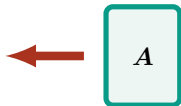
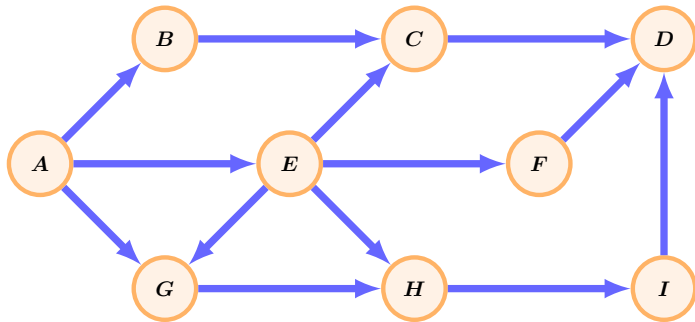
### DEPTH-FIRST-SEARCH :

1. Select root vertex and add to a stack
2. **While** stack is not-empty:
  3. Inspect the top element  $u$  on the stack.
  4. Add any unvisited neighbours of  $u$  to the stack.
  5. Remove  $u$  from the stack.

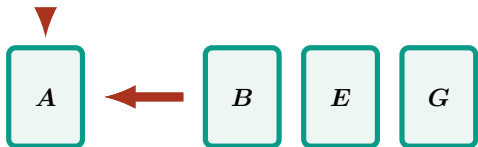
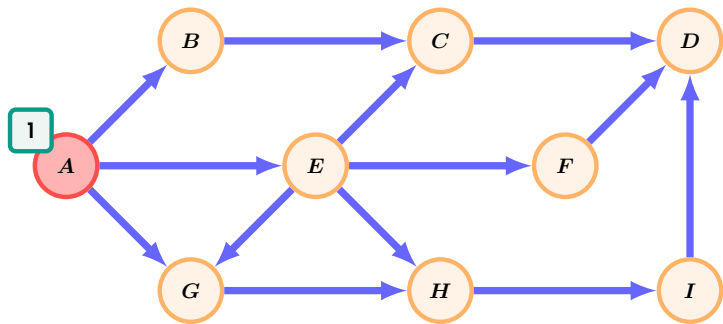
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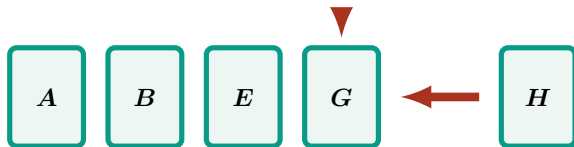
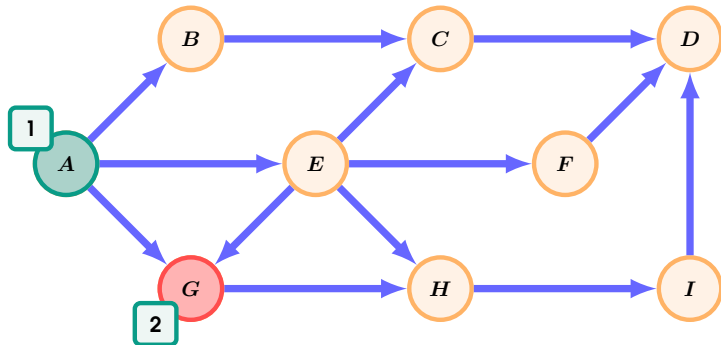
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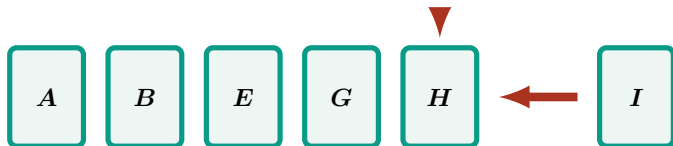
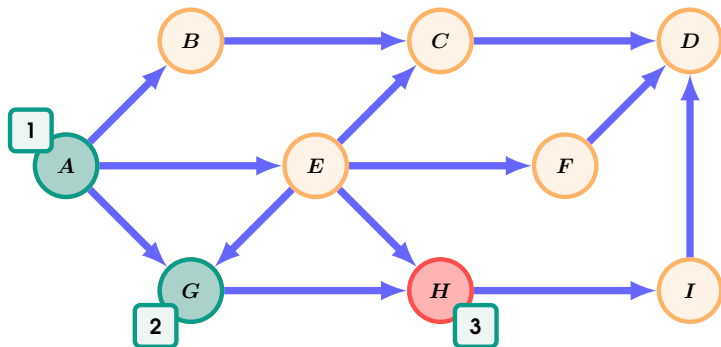
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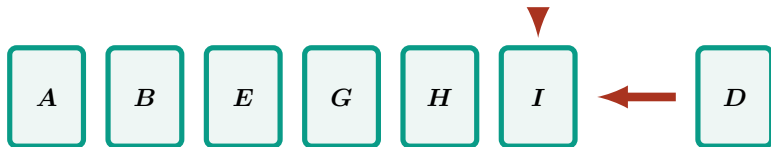
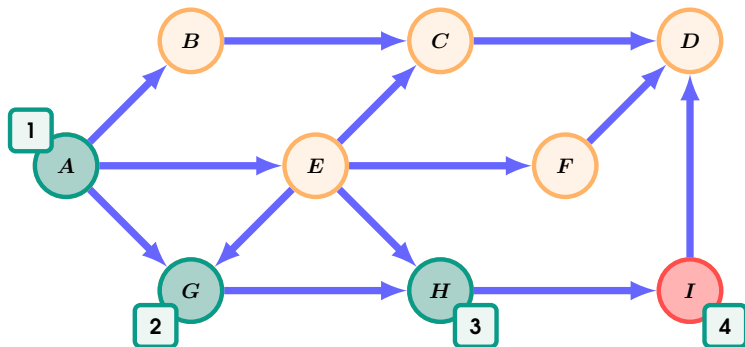
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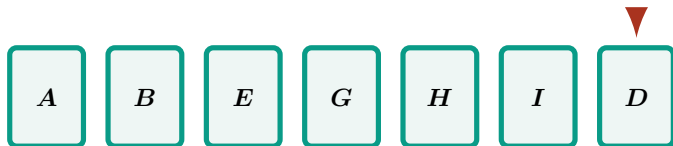
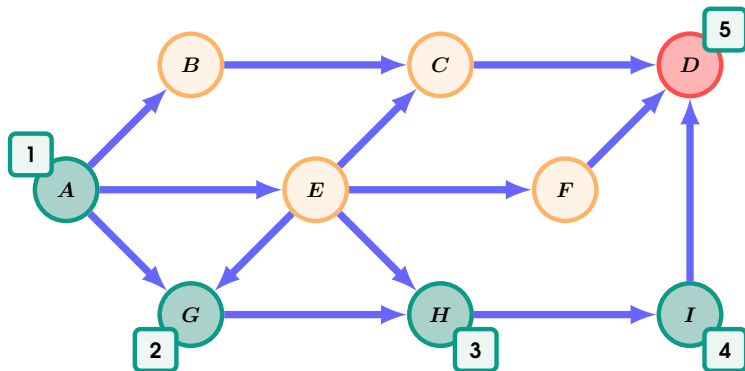
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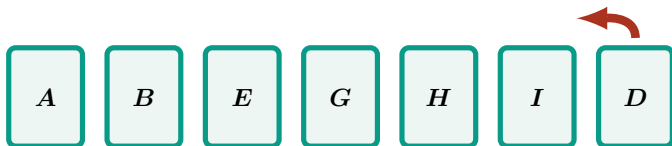
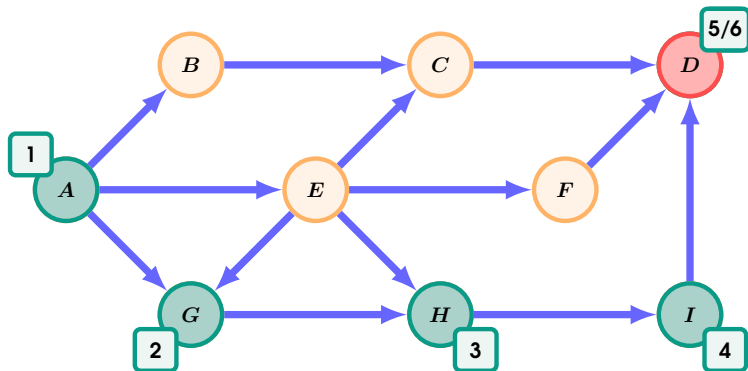
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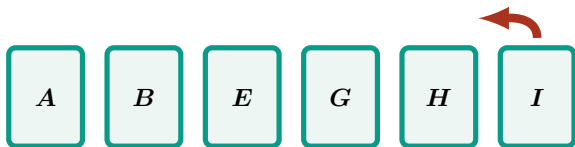
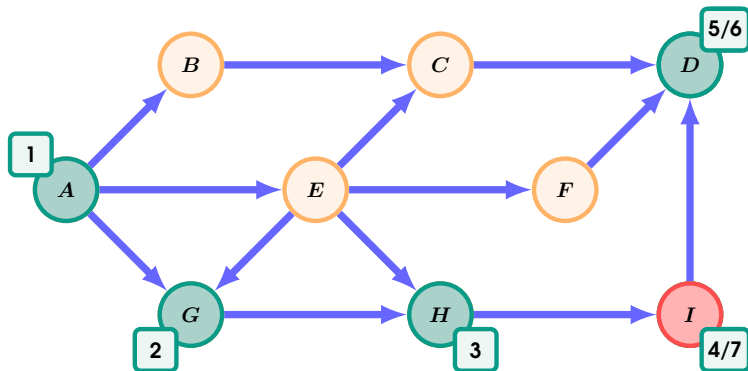
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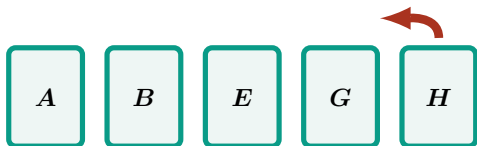
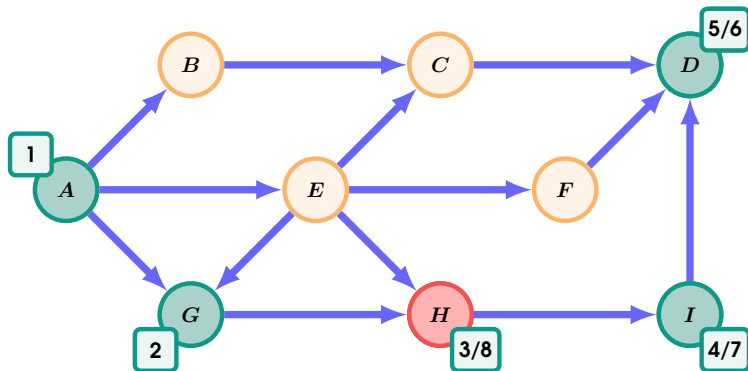
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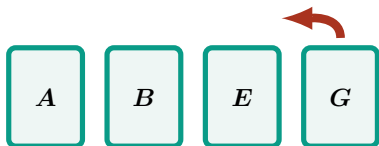
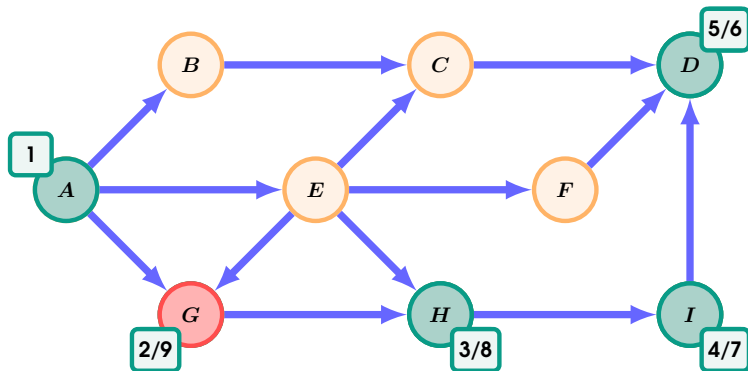
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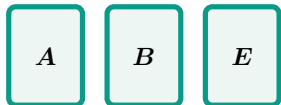
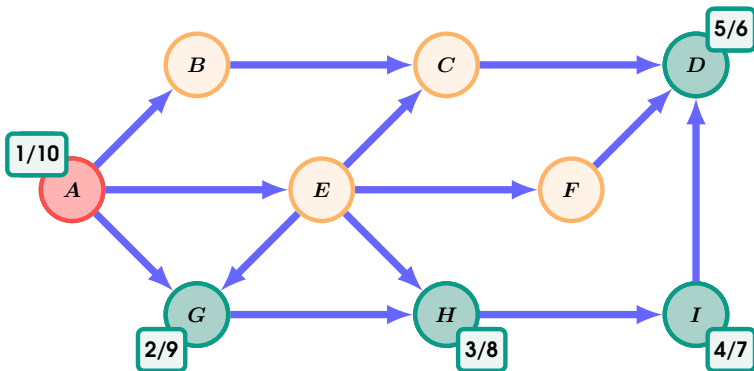
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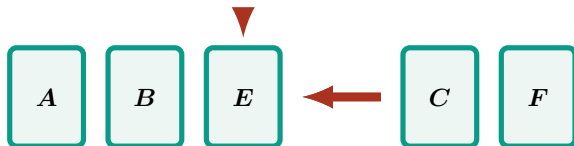
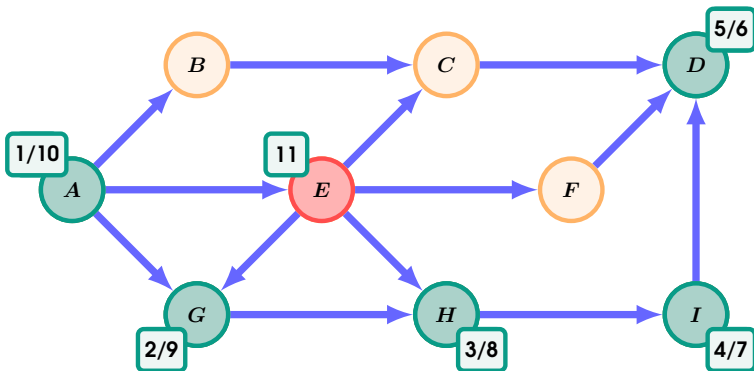
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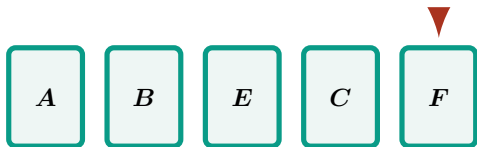
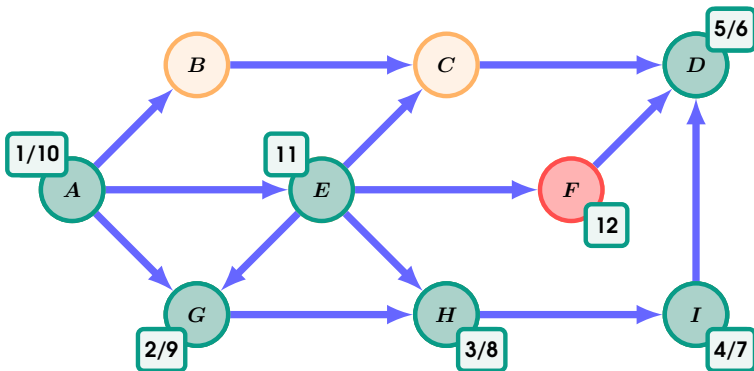
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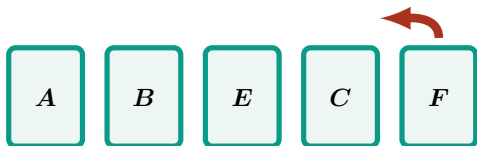
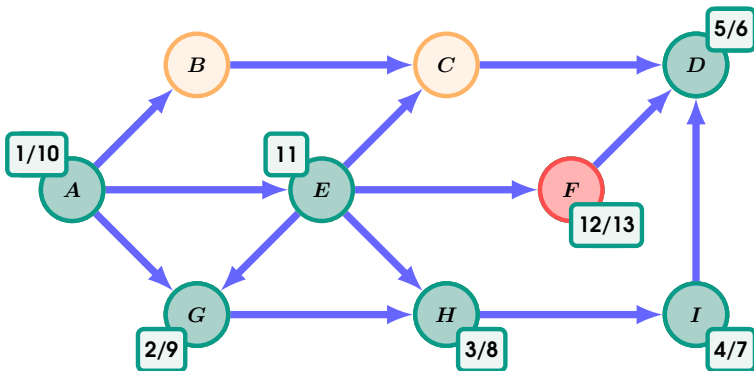
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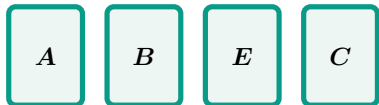
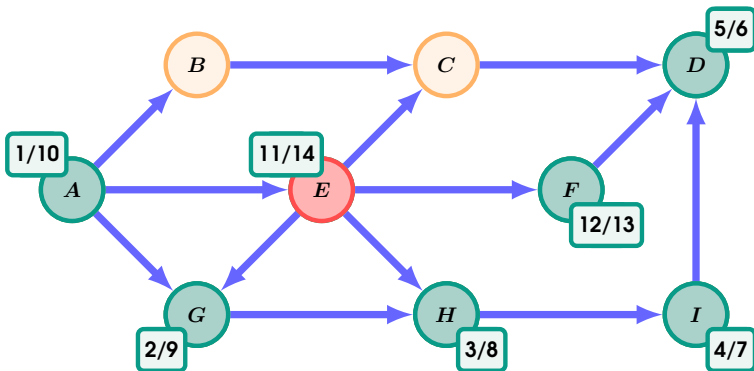
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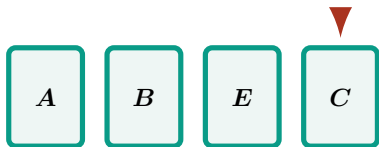
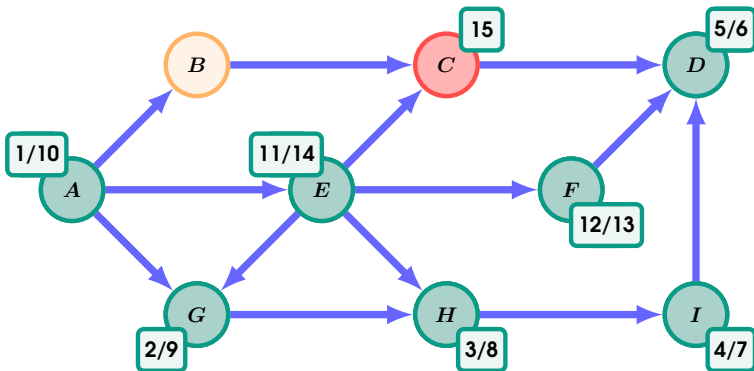
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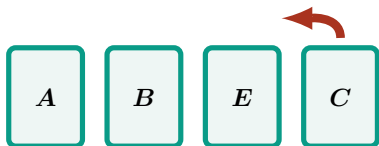
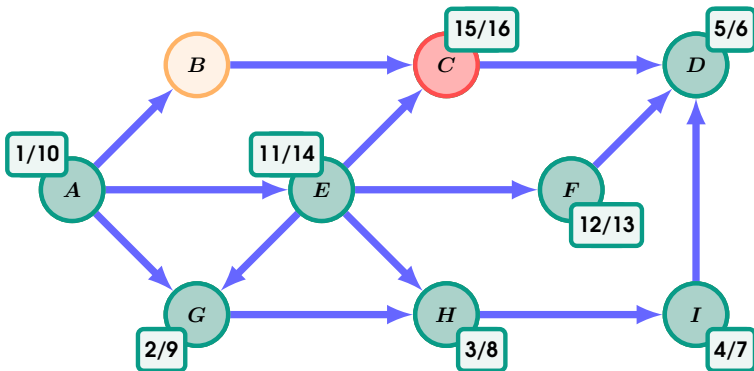
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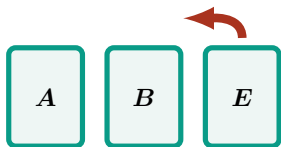
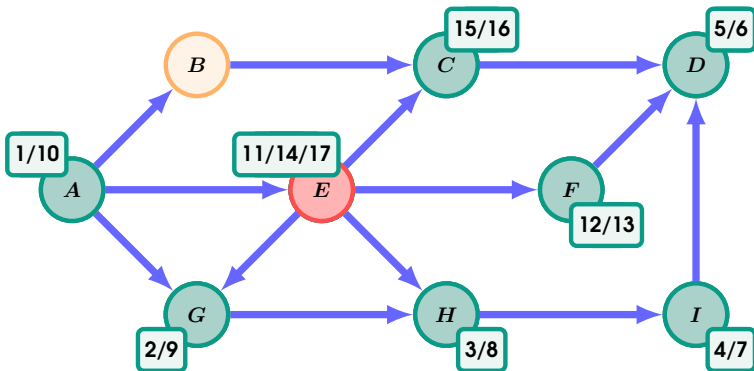
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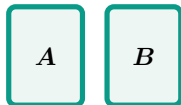
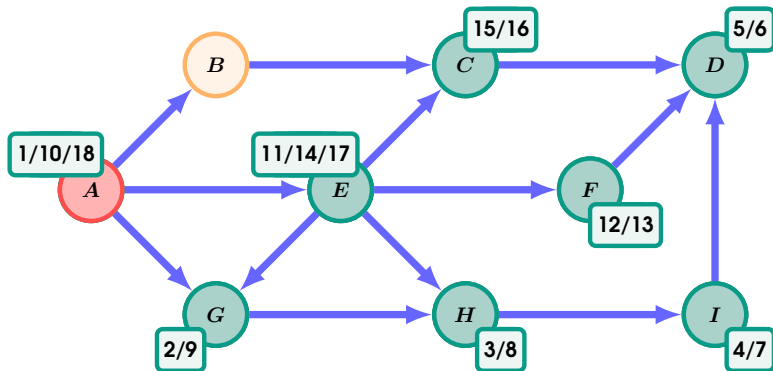
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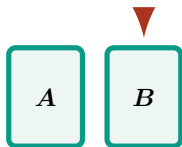
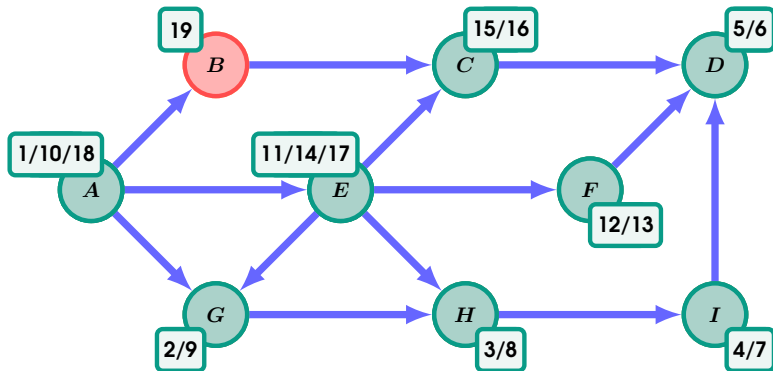
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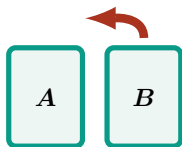
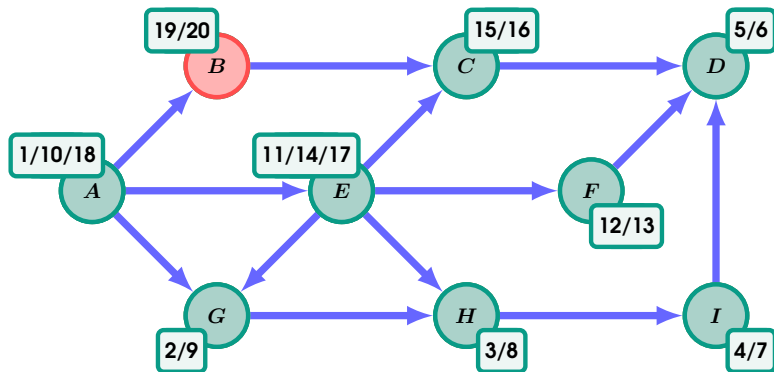
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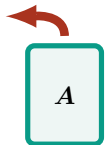
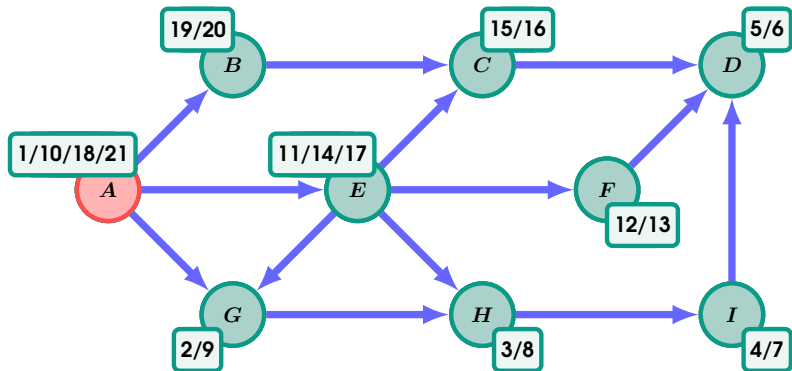
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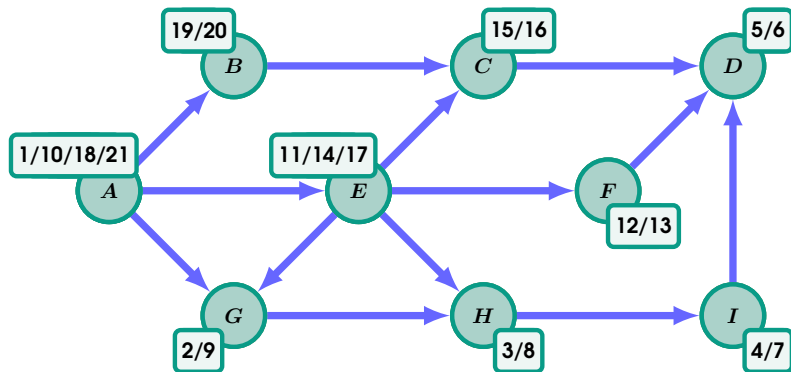
## Depth-First-Search



## Depth-First-Search



## Depth-First-Search



**Stack Empty so Terminate!**

## Depth-First-Search

**Theorem** The termination time for Depth-First-Search is  $O(|V| + |E|)$ .

### *Proof:*

- The proof is almost identical to that of the Breadth-First-Search, just replacing the **queue** for a **stack**.

**Q.E.D.**

# End of Slides!

