

5CCS2FC2: Foundations of Computing II

Introduction to Graph Algorithms

Week 4

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Topological Sorting

Directed Acyclic Graphs (DAGs)

- Directed Acyclic Graphs (DAGs)

- A graph $G = (V, E)$ is said to be a **directed acyclic graph** if
 - it is **irreflexive**, and
 - does not contain any **cycles** of length ≥ 2

if there is a path $u \rightsquigarrow v$ then there is no path $v \rightsquigarrow u$

Example: Software Dependency Graphs

<http://cxf.apache.org/docs/cxf-dependency-graphs.html>

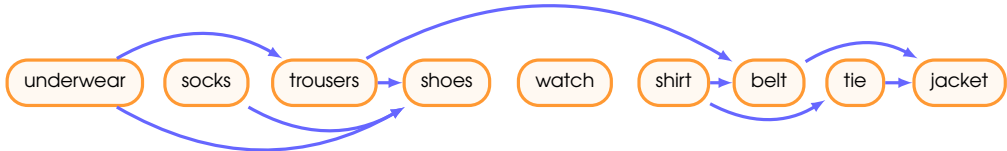
Topological Sorting

- Topological Sort

- A **topological sorting** of a Directed Acyclic Graph $G = (V, E)$ is an ordered list of vertices $L = \langle v_1, v_2, \dots, v_n \rangle$ such that

If $v_i \rightsquigarrow v_j$ then $i < j$ for all $i, j \leq n$

(i.e. all the arrows point 'downstream' from v_1 to v_n)

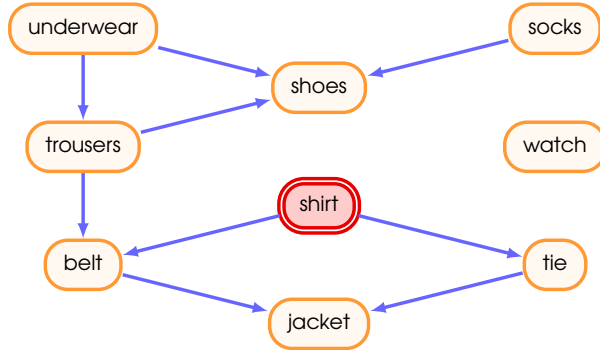


Topological Sorting

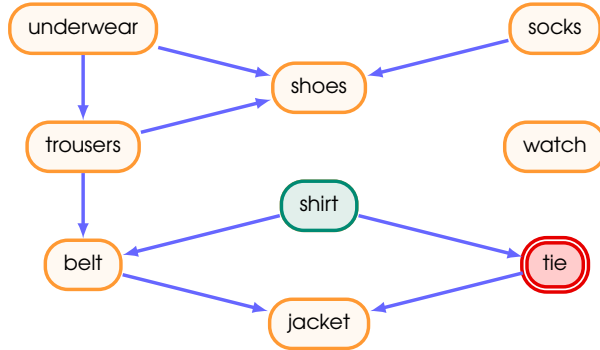
TOPOLOGICAL-SORT (G) :

1. **While** vertices remain unsorted:
 2. Select any unsorted vertex and add to a stack
 3. **While** stack is not-empty:
 4. Inspect the top element u on the stack.
 5. **If** u has no unsorted neighbours:
 6. Pop u from the stack and add to list L .
 7. **Else:**
 8. Add any unsorted neighbours of u to the stack.
9. **Return** sorted list L

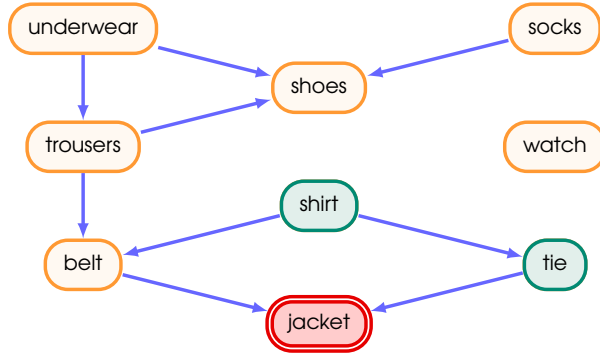
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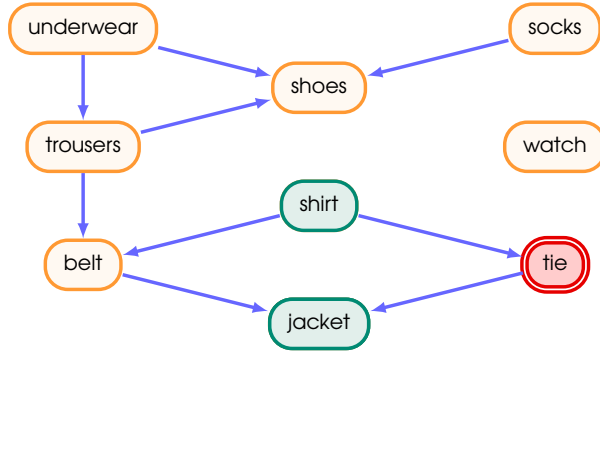
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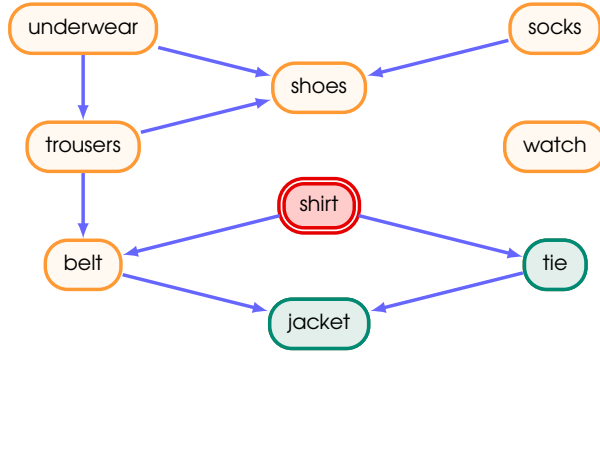
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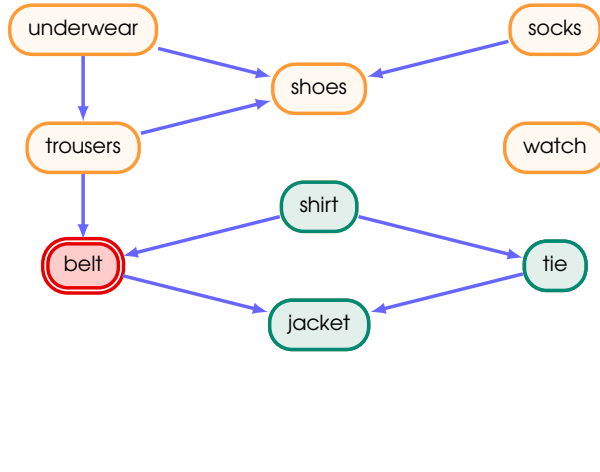
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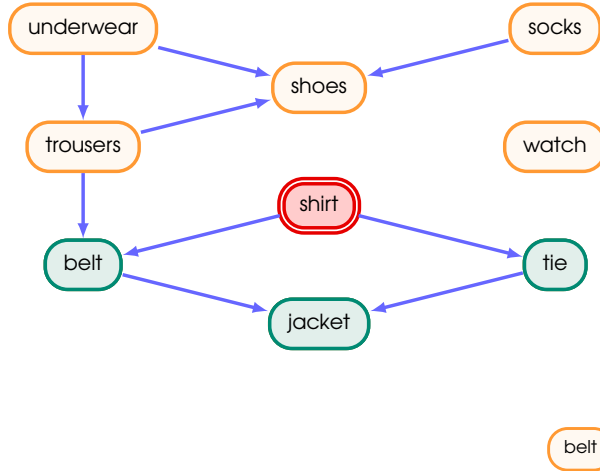
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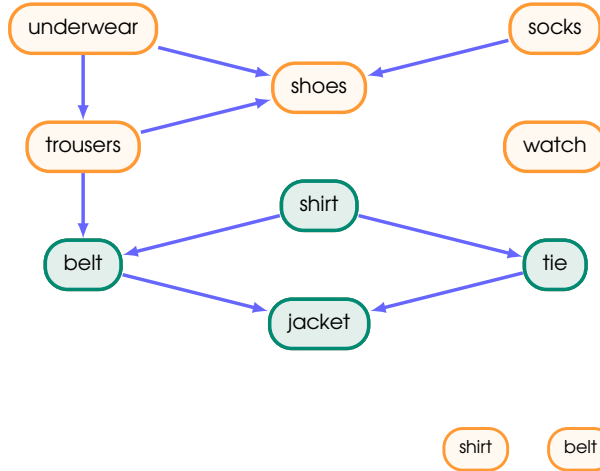
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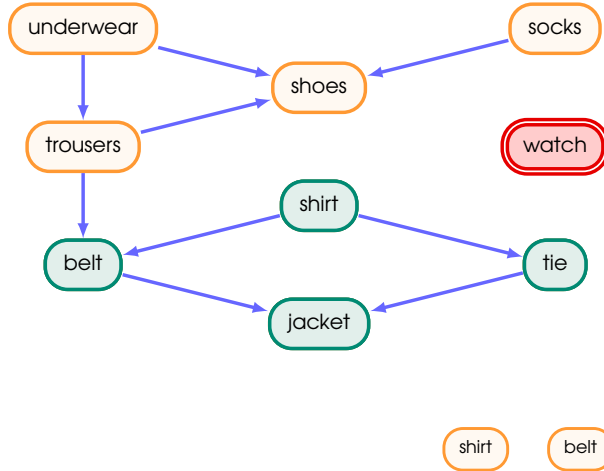
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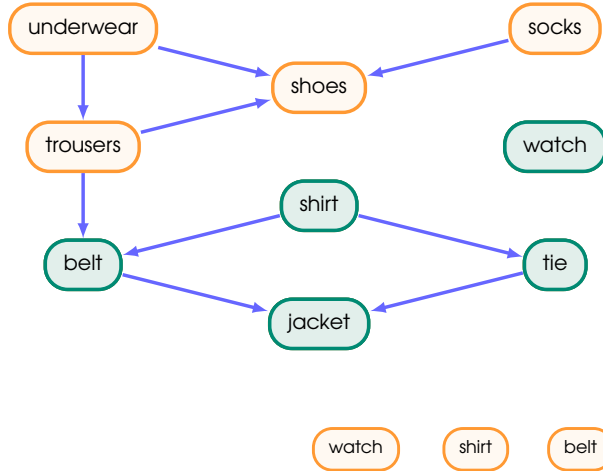
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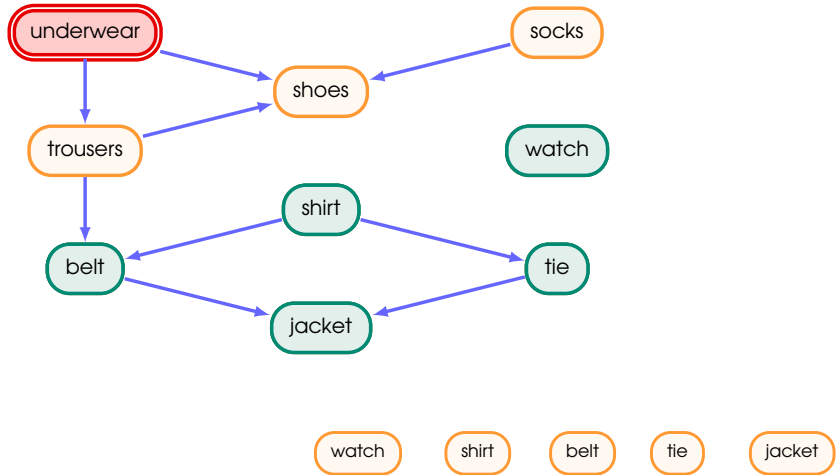
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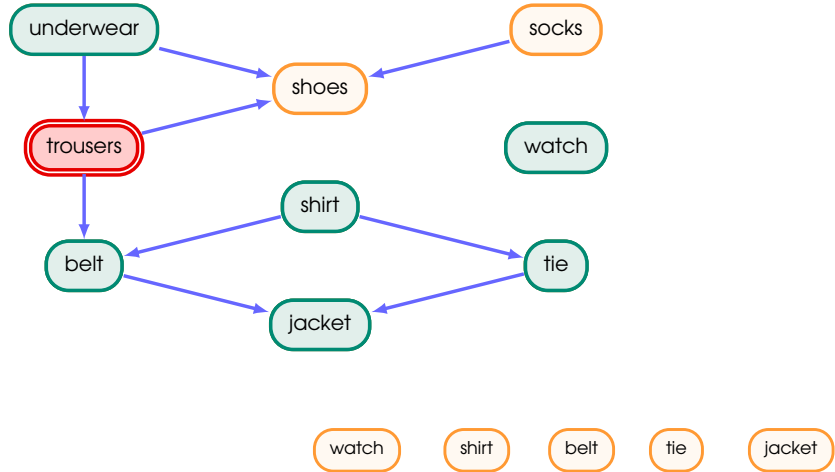
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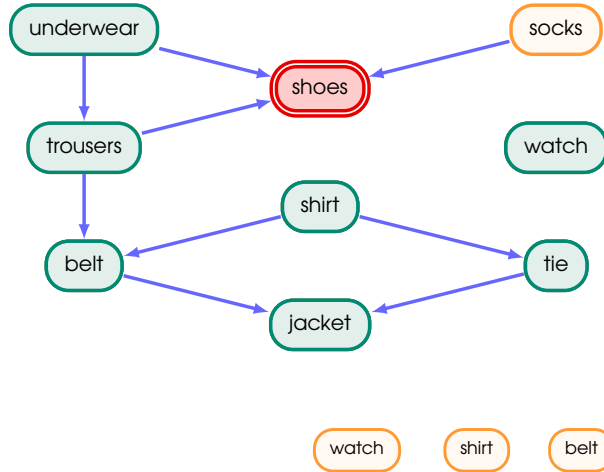
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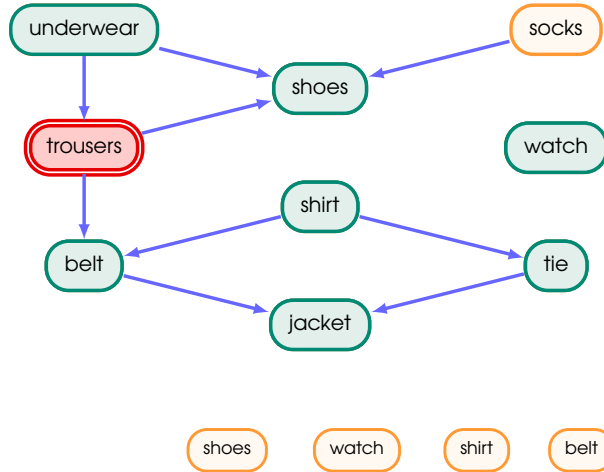
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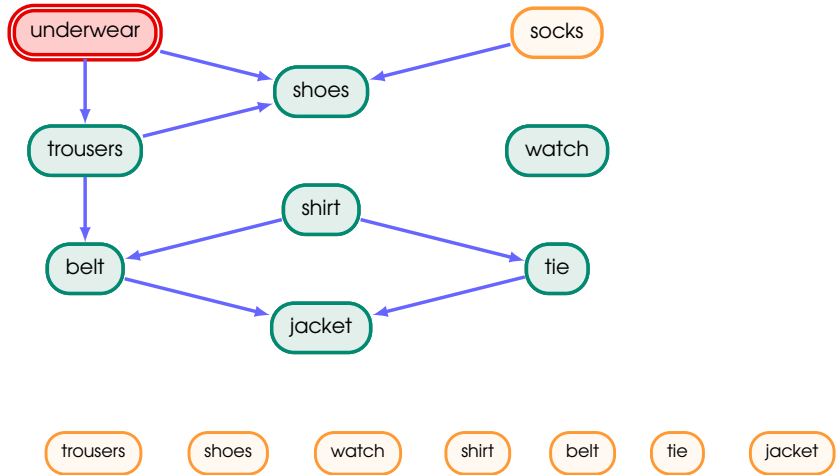
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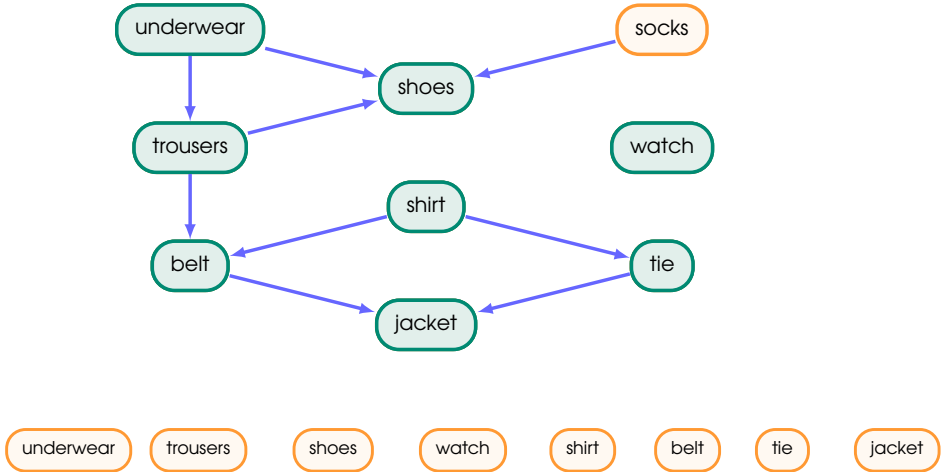
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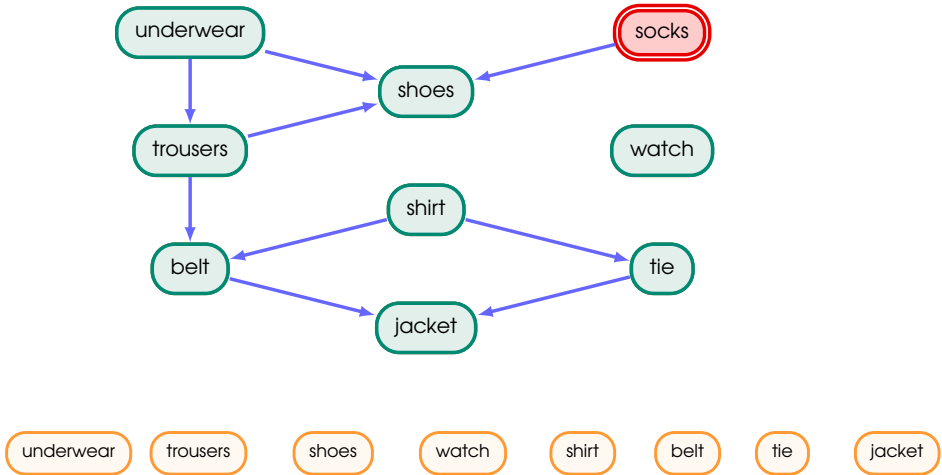
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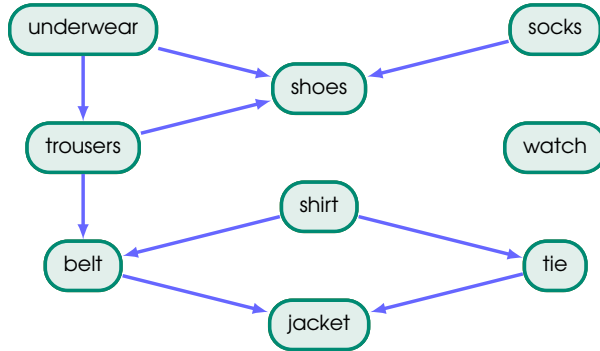
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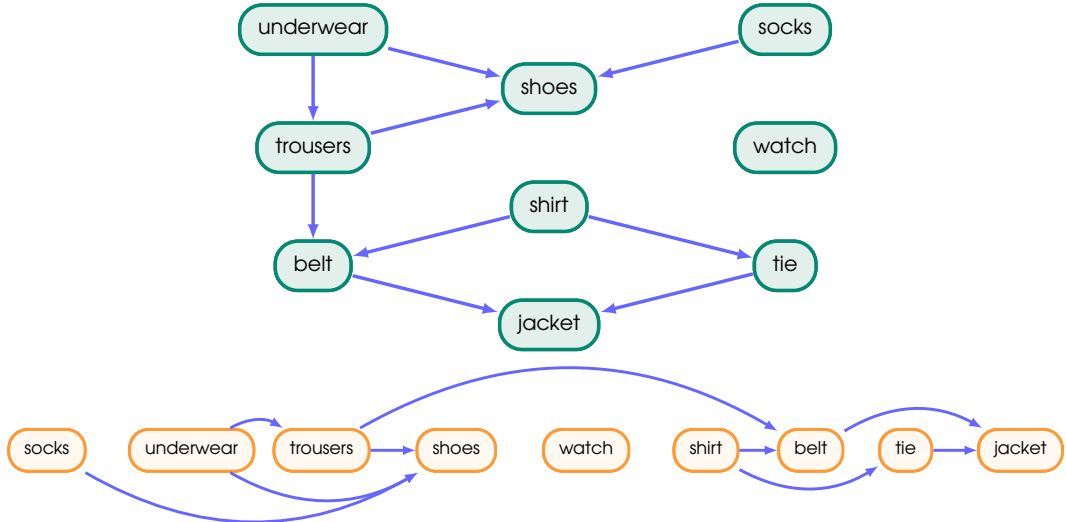
Topological Sorting



Topological Sorting



Topological Sorting



Strongly Connected Components

Some Definitions

- Strong Connected Component (SCC)

- We can define an **binary relation** $\sim \subseteq V \times V$ on the set of vertices such that

$$u \sim v \iff u = v \text{ or there is a path } u \rightsquigarrow v \text{ and a path } v \rightsquigarrow u$$

- A **strongly connected component** is a *maximal* subset $C \subseteq V$, such that

$$u \sim v \quad \text{for all } u, v \in C$$

$$[u]_{\equiv} = \{v \in V : v \sim u\}$$

Some Definitions

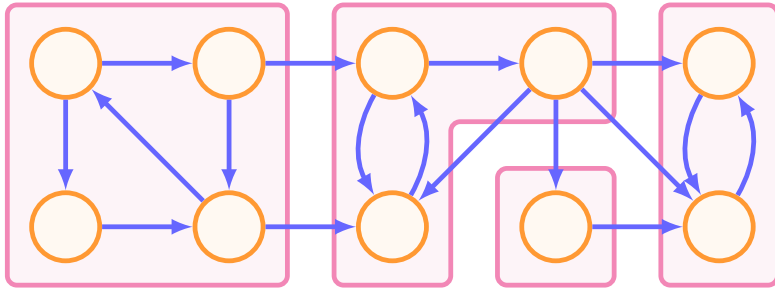
- **Component Graph**

- The **component graph** of G is a new graph $G^{\text{scc}} = (V^{\text{scc}}, E^{\text{scc}})$, where

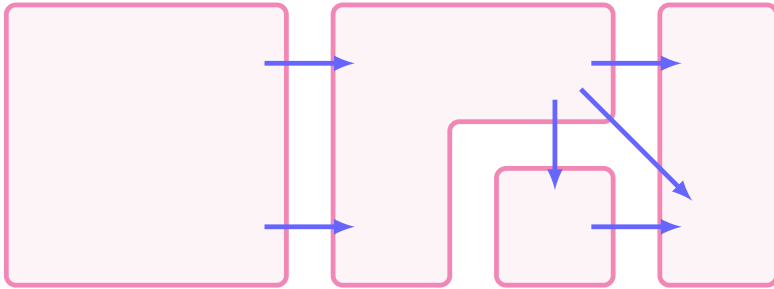
$$V^{\text{scc}} = \{ \text{all strongly connected components of } G \}$$

$$(A, B) \in E^{\text{scc}} \iff (a, b) \in E \quad \text{for some } a \in A \text{ and } b \in B, \\ \text{where } A \neq B$$

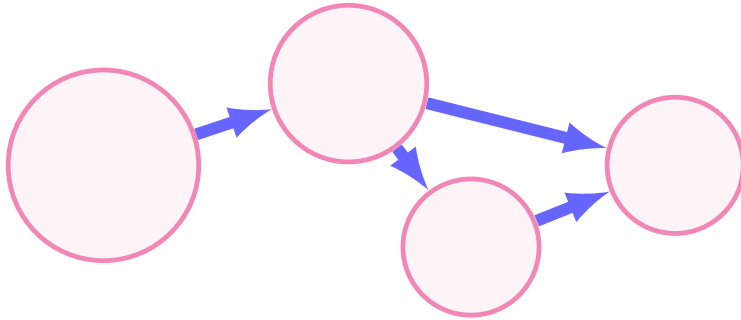
Some Definitions



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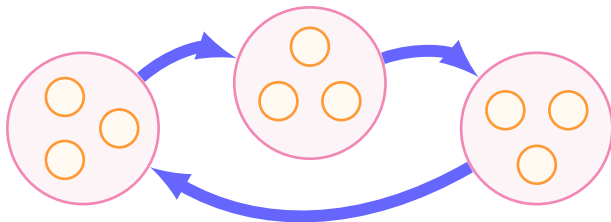


Strongly Connected Components

Theorem The component graph G^{SCC} is a directed acyclic graph, for any directed graph G .

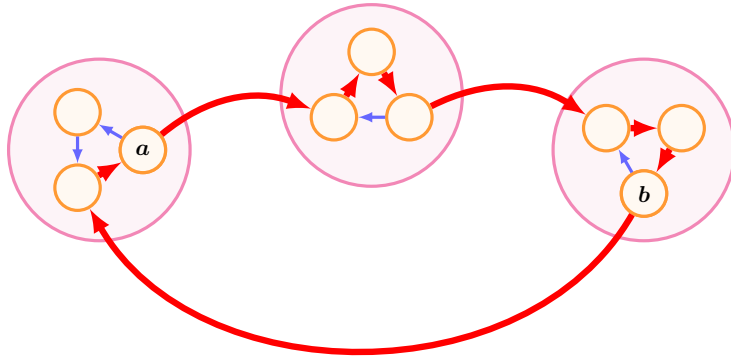
Proof:

Step 1) Suppose, for **contradiction**, that there is a cycle $\langle C_1, C_2, \dots, C_n, C_1 \rangle$:



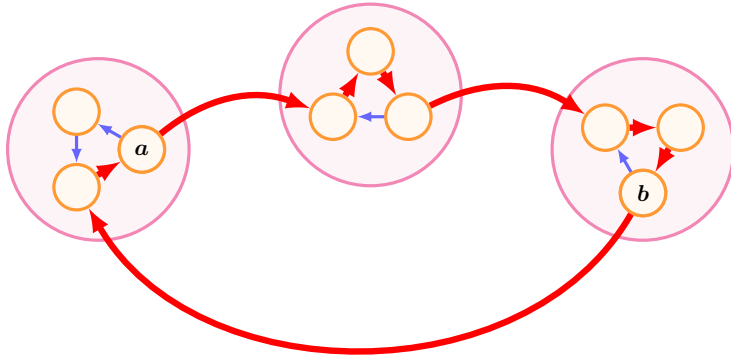
Strongly Connected Components

Step 2) Choose any $a \in C_i$ and $b \in C_j$ for some $C_i \neq C_j$.



Strongly Connected Components

Step 3) It follows that there must be a path $a \rightsquigarrow b$ and $b \rightsquigarrow a$



Strongly Connected Components

Step 4) Hence, by **contradiction**, G^{SCC} cannot contain any cycles, as required.

Q.E.D.

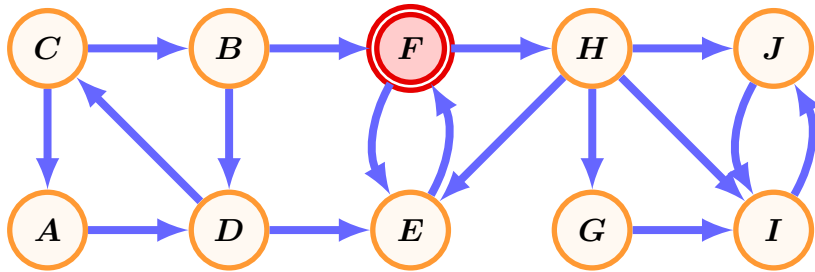
Finding Strongly Connected Components

Strongly Connected Components

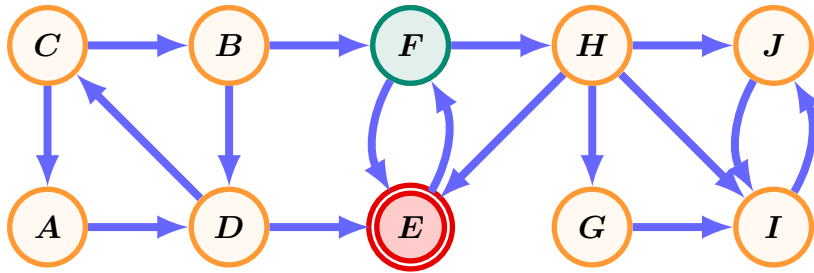
SCC (G):

1. Call TOPOLOGICAL-SORT (G) to generate ordered list L
2. **While** L is not-empty:
 3. Inspect the first element u in the list.
 4. Perform DFS on the *transpose graph* G^T from u
 5. Add all visited vertices to a new component S
 6. Remove these vertices from the list L .
7. **Return** all components identified

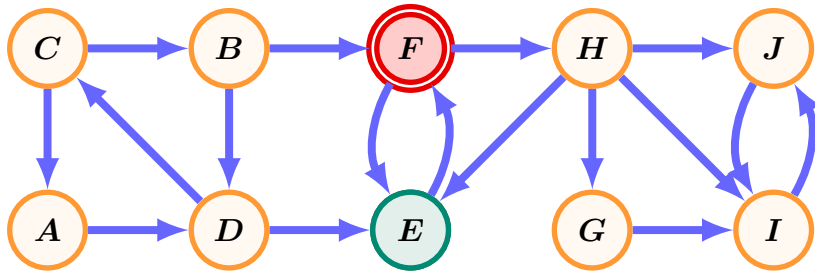
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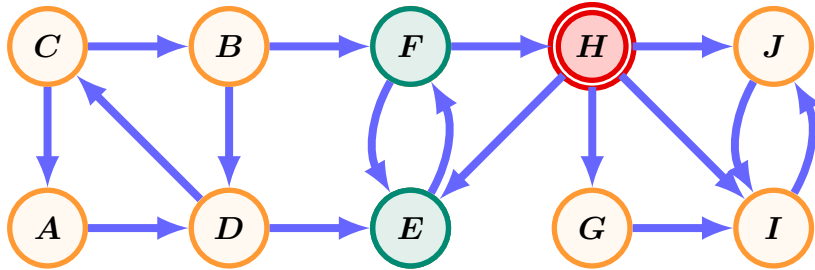
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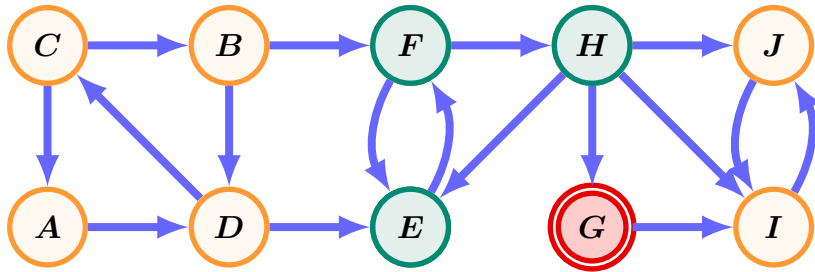
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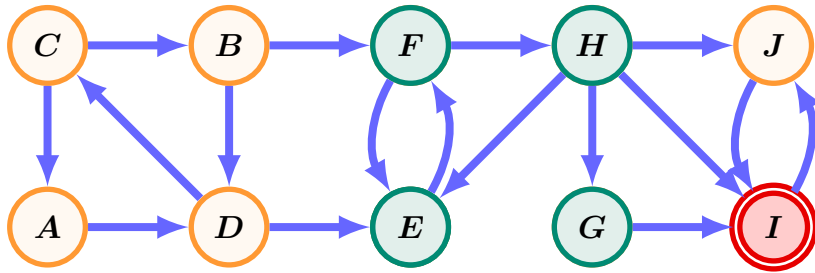
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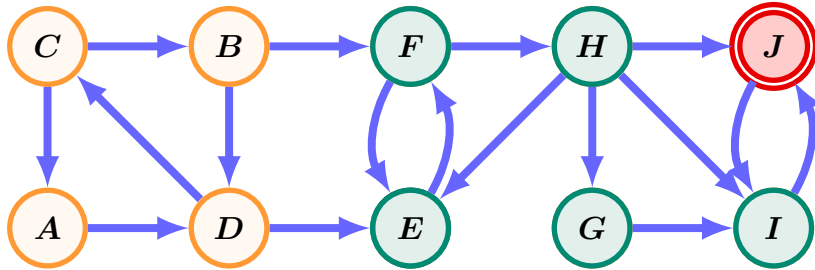
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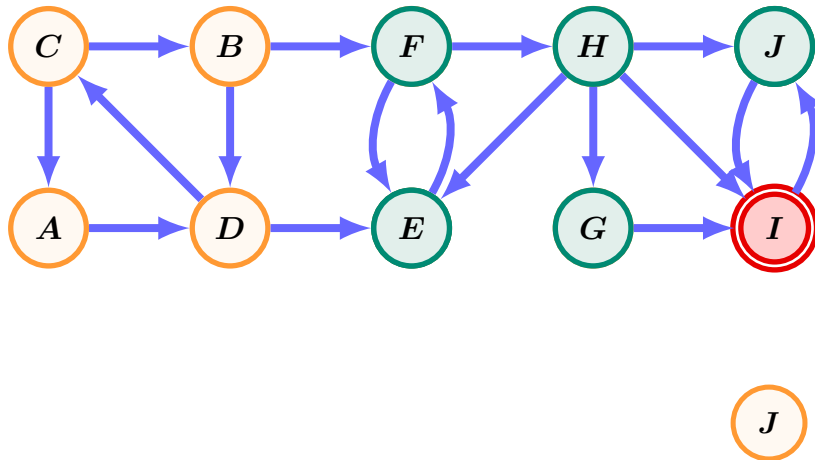
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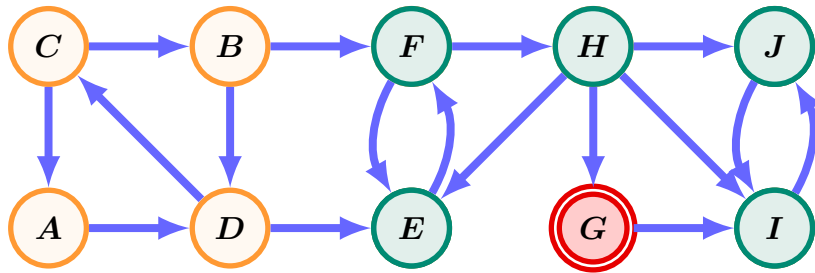
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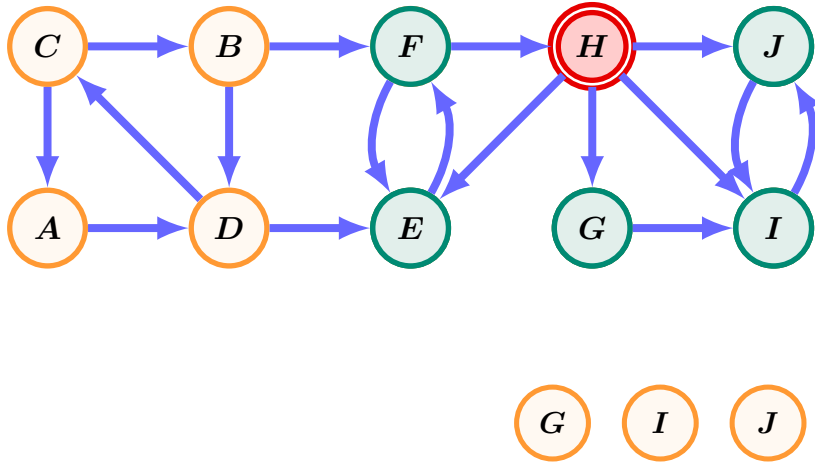
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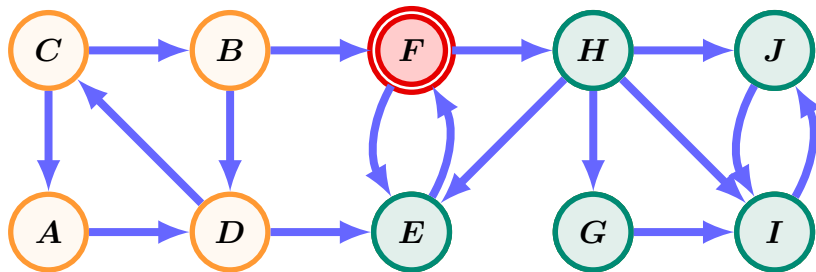
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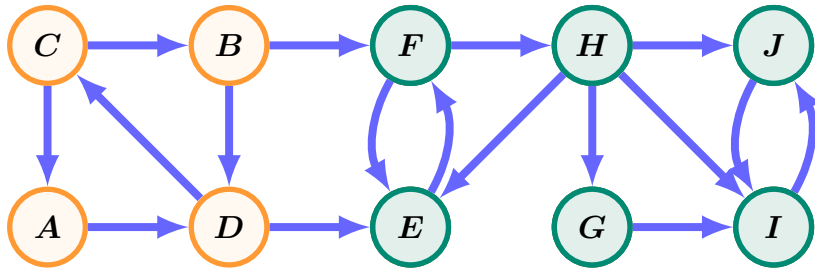
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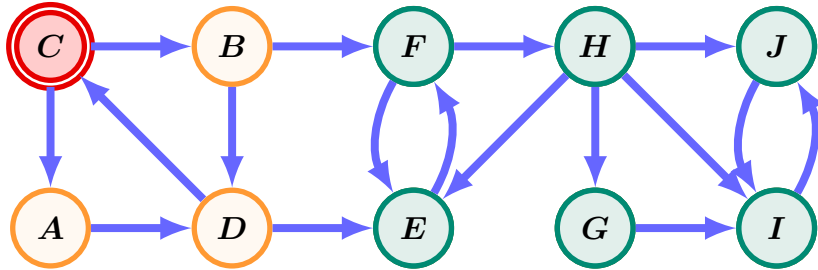
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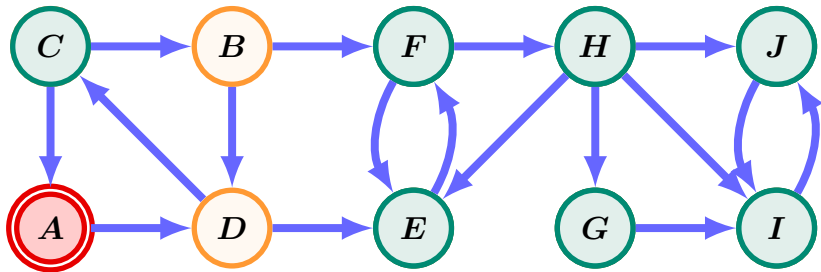
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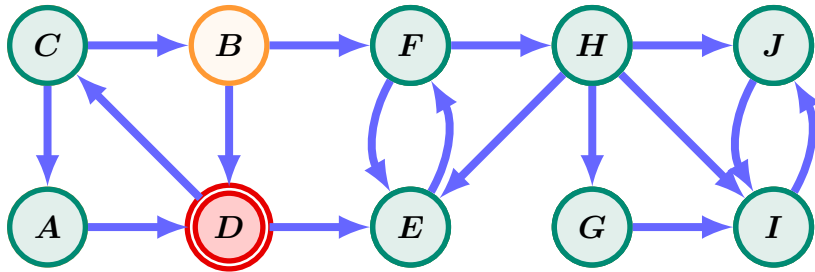
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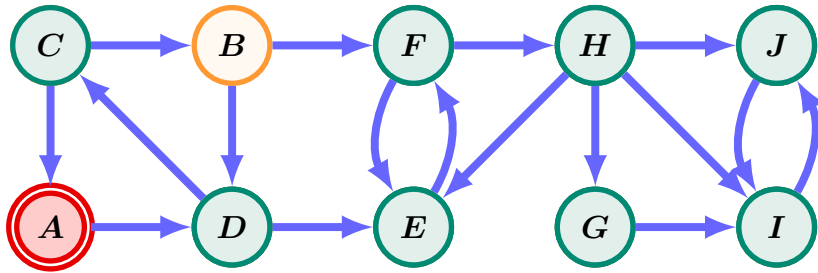
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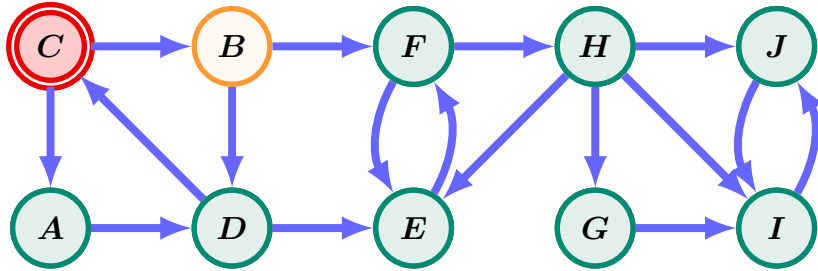
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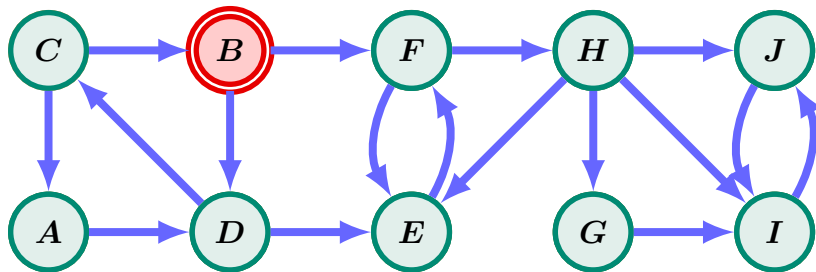
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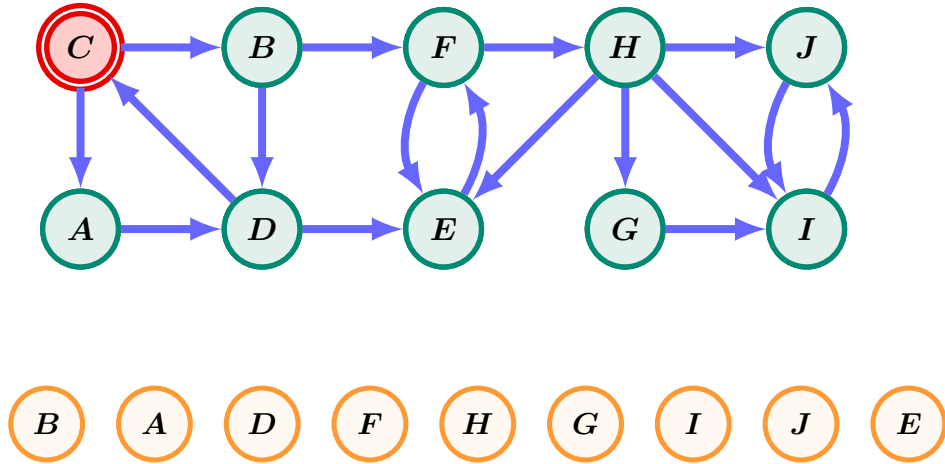
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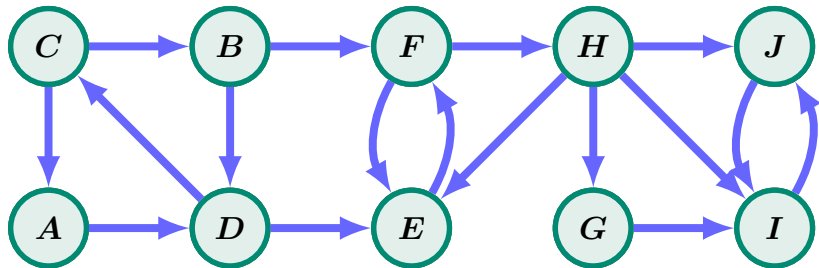
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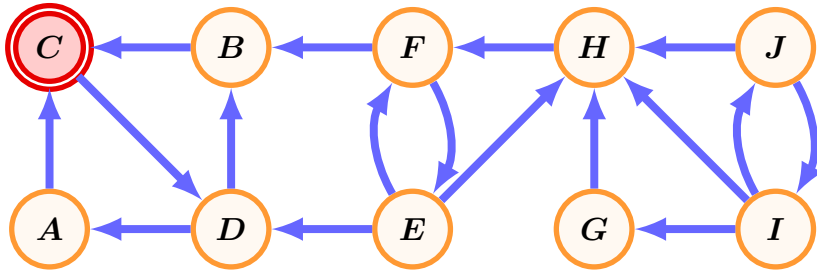
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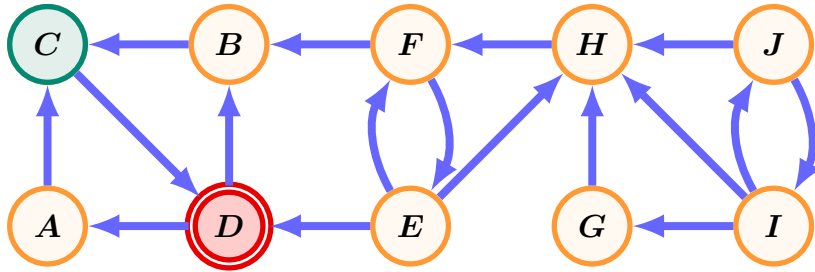
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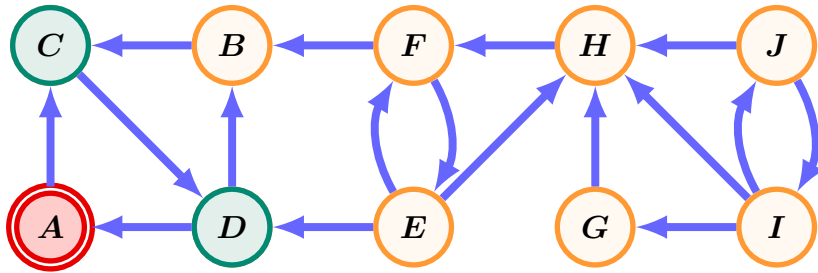
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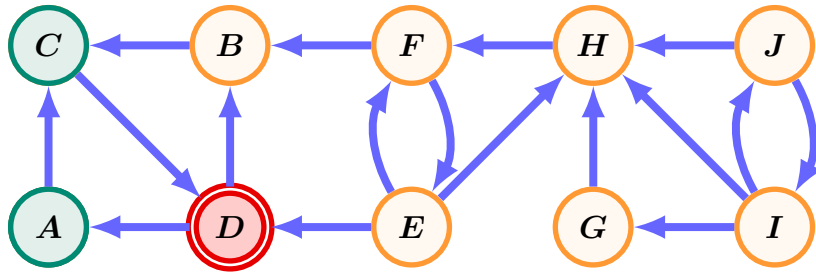
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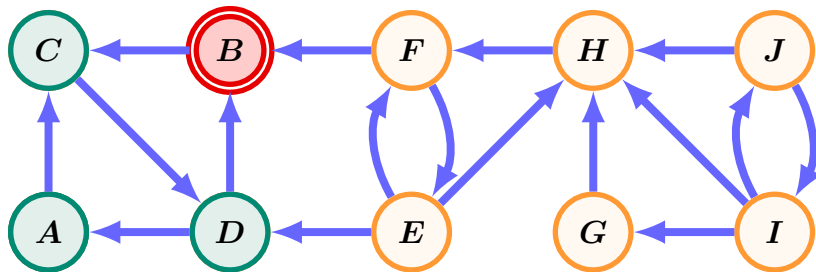
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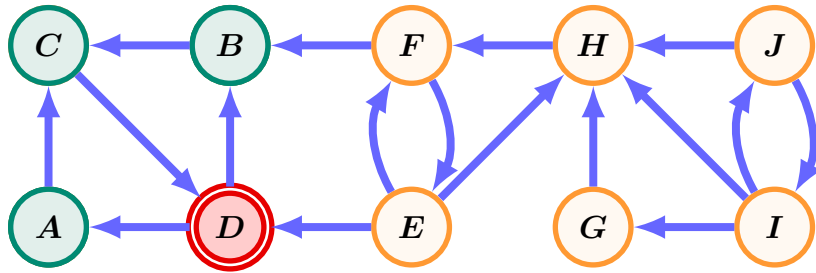
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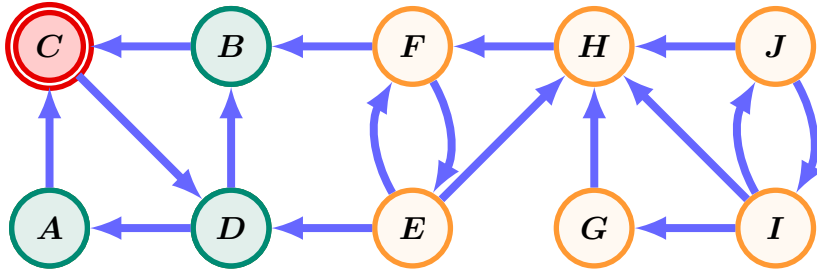
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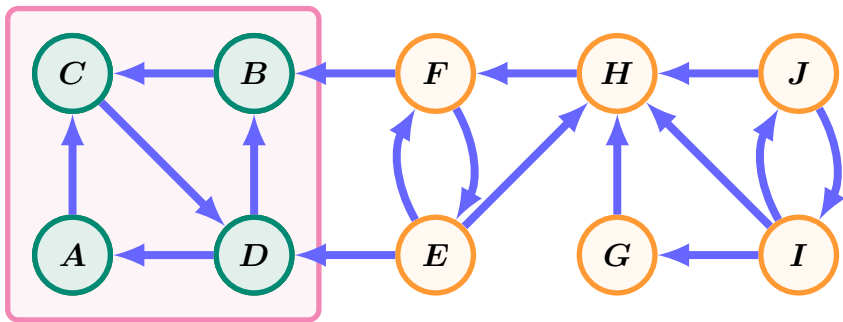
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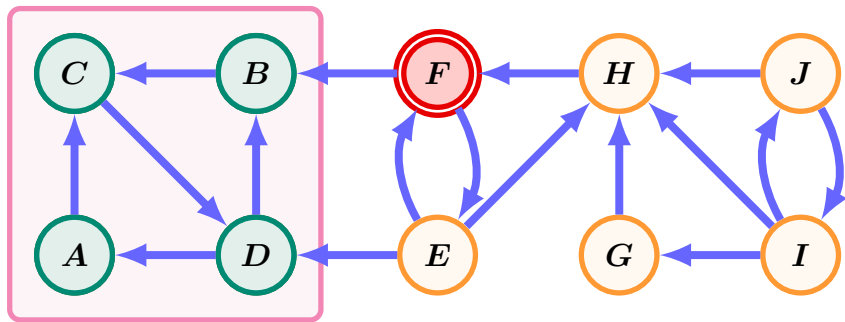
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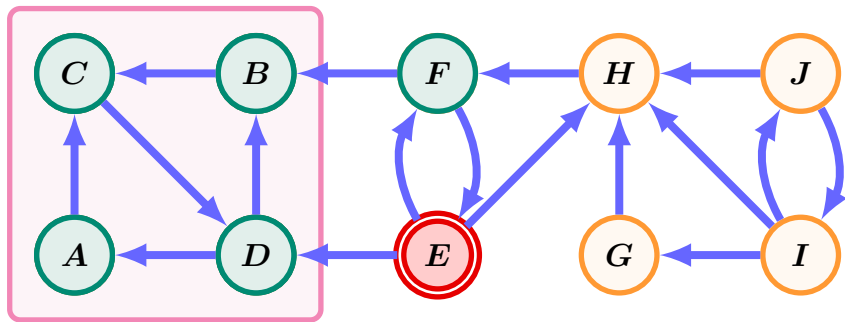
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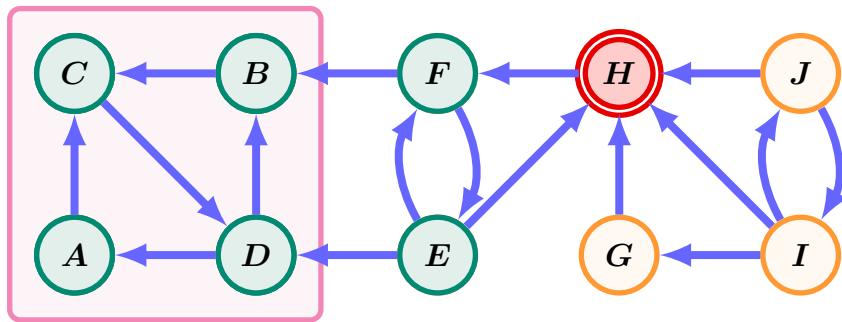
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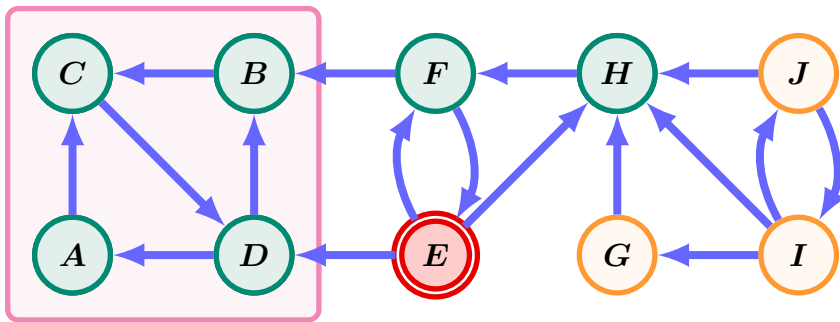
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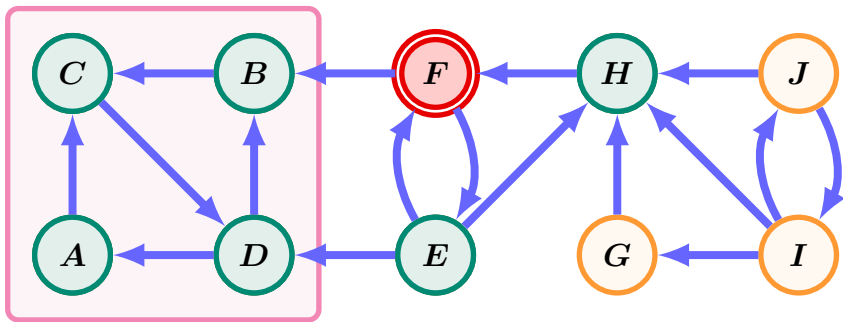
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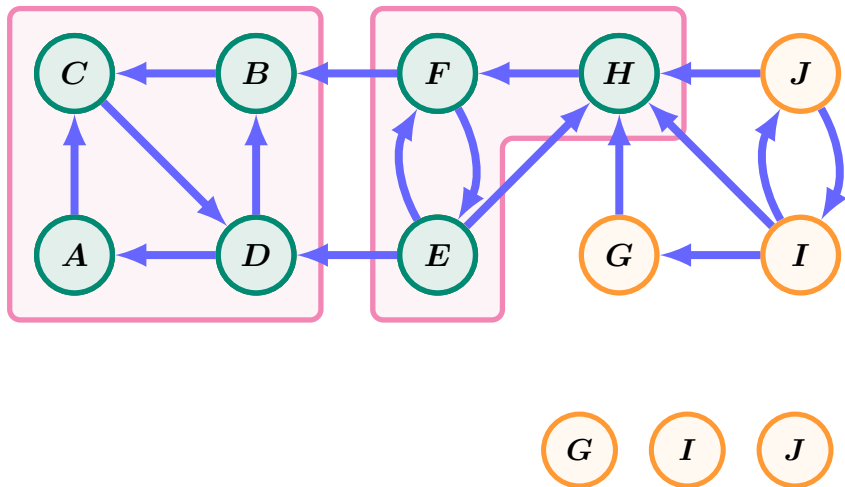
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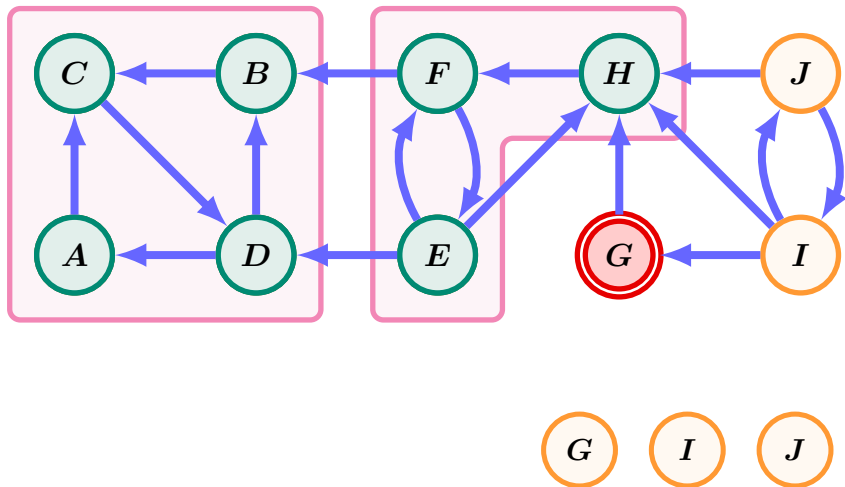
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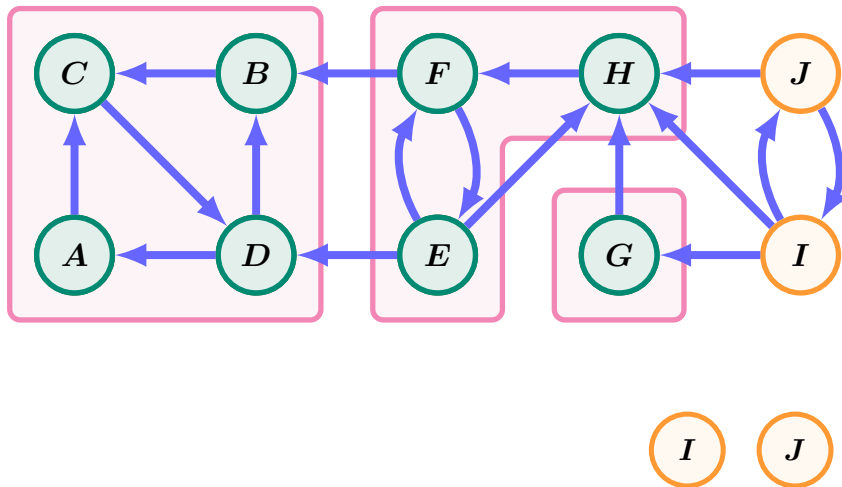
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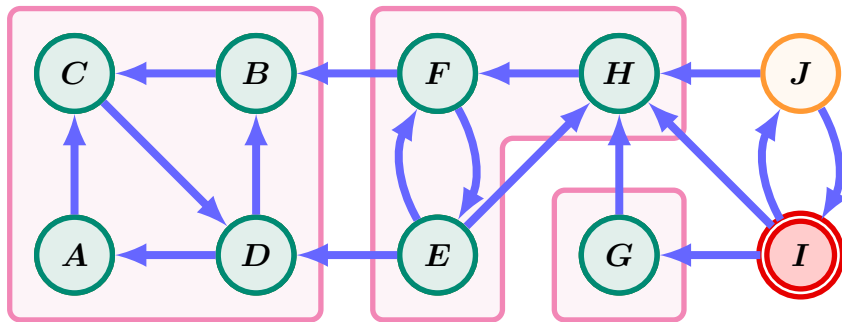
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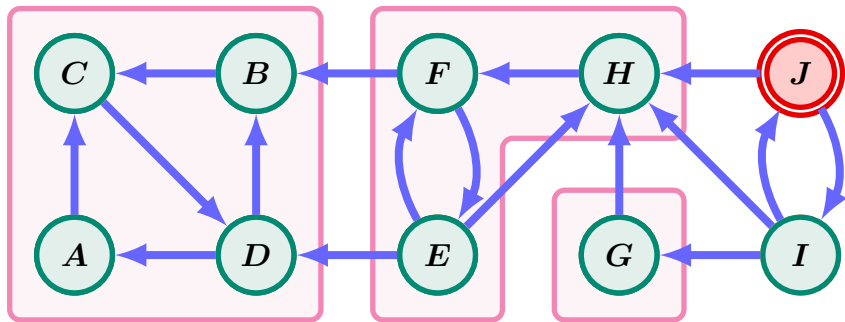
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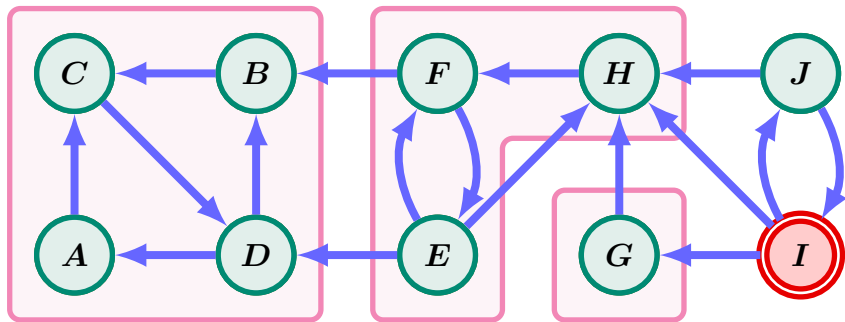
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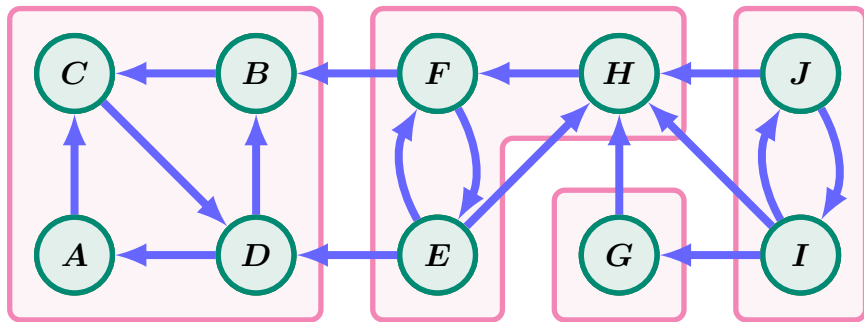
Strongly Connected Components



Strongly Connected Components



Strongly Connected Components



End of Slides!

